## MATH 233: ASSIGNMENT 3

DUE: FRIDAY, FEBRUARY 16, 11:59PM ON GRADESCOPE
UNIVERSITY OF ROCHESTER, SPRING 2024

Problem 1 (6.6.1). The ciphertext YIFZMA was encrypted by a Hill cipher with matrix

$$
\left(\begin{array}{cc}
9 & 13 \\
2 & 3
\end{array}\right) \quad(\bmod 26)
$$

Find the plaintext.

Solution. Let $A$ be the given matrix $(\bmod 26)$. Then

$$
\operatorname{det} A \equiv 9 \cdot 3-13 \cdot 2 \equiv 1 \quad(\bmod 26),
$$

so

$$
A^{-1} \equiv\left(\begin{array}{cc}
3 & -13 \\
-2 & 9
\end{array}\right) \quad(\bmod 26)
$$

Then

$$
\begin{aligned}
& \text { YI } \quad \rightarrow \quad\left(\begin{array}{ll}
24 & 8
\end{array}\right) A^{-1} \equiv\left(\begin{array}{ll}
56 & -240
\end{array}\right) \equiv\left(\begin{array}{ll}
4 & 20
\end{array}\right) \quad \rightarrow \quad \text { eu } \\
& \text { FZ } \quad \rightarrow \quad\left(\begin{array}{ll}
5 & 25
\end{array}\right) A^{-1} \equiv\left(\begin{array}{ll}
-35 & 160
\end{array}\right) \equiv\left(\begin{array}{ll}
17 & 4
\end{array}\right) \quad \rightarrow \quad \text { re } \\
& \text { MA } \quad \rightarrow \quad\left(\begin{array}{ll}
12 & 0
\end{array}\right) A^{-1} \equiv\left(\begin{array}{ll}
36 & -156
\end{array}\right) \equiv\left(\begin{array}{ll}
10 & 0
\end{array}\right) \quad \rightarrow \quad \text { ka }
\end{aligned}
$$

shows that the plaintext is eureka.
Problem 2 (6.6.5). Eve captures Bob's Hill cipher machine, which uses a $2 \times 2$ matrix $M$ mod 26. She tries a chosen plaintext attack. She finds that the plaintext ba encrypts to HC and the plaintext zz encrypts to GT. What is the matrix $M$ ?

Solution. From the given information, we construct a matrix equation as

$$
\left(\begin{array}{cc}
1 & 0 \\
-1 & -1
\end{array}\right) M \equiv\left(\begin{array}{cc}
7 & 2 \\
6 & 19
\end{array}\right) \quad(\bmod 26) .
$$

(We used -1 instead of $z=25$ for simplicity.) The first matrix is invertible mod 26, with inverse

$$
\left(\begin{array}{cc}
1 & 0 \\
-1 & -1
\end{array}\right)
$$

(note that it is the same as the original matrix), so

$$
M \equiv\left(\begin{array}{cc}
1 & 0 \\
-1 & -1
\end{array}\right)\left(\begin{array}{cc}
7 & 2 \\
6 & 19
\end{array}\right) \equiv\left(\begin{array}{cc}
7 & 2 \\
-13 & -21
\end{array}\right) \equiv\left(\begin{array}{cc}
7 & 2 \\
13 & 5
\end{array}\right) \quad(\bmod 26) .
$$

Note. Problems 3 and 4 are about designing the known plaintext attack for certain cryptosystems. For each problem, assume that you can use only one plaintext. You should choose the plaintext wisely, so that you are able to find the entire key(s) regardless of the ciphertext that you get. Furthermore, you should try to keep the length of your plaintext as short as possible. A full credit may not be given if your plaintext is too long.

Problem 3 (6.6.4). Consider the following combination of Hill and Vigenère ciphers: The key consists of three $2 \times 2$ matrices, $M_{1}, M_{2}, M_{3}(\bmod 26)$. The plaintext letters are represented as integers mod 26. The first two are encrypted by $M_{1}$, the next two by $M_{2}$, the 5 th and 6 th by $M_{3}$. This is repeated cyclically, as in the Vigenère cipher. Explain how to do a chosen plaintext attack on this system. Assume that you know that three $2 \times 2$ matrices are being used. State explicitly what plaintext you would use and how you would use the outputs.

Solution. Use bababaababab $=101010010101$ as the plaintext, and suppose that the ciphertext is $c_{1} c_{2} \ldots c_{12}$ (where each $c_{i}$ is represented as an integer mod 26). Then we have

$$
\left(\begin{array}{ll}
1 & 0 \\
0 & 1
\end{array}\right) M_{1} \equiv\left(\begin{array}{ll}
c_{1} & c_{2} \\
c_{7} & c_{8}
\end{array}\right) \quad \Rightarrow \quad M_{1} \equiv\left(\begin{array}{ll}
c_{1} & c_{2} \\
c_{7} & c_{8}
\end{array}\right) \quad(\bmod 26)
$$

and similarly

$$
M_{2} \equiv\left(\begin{array}{cc}
c_{3} & c_{4} \\
c_{9} & c_{10}
\end{array}\right) \quad \text { and } \quad M_{3} \equiv\left(\begin{array}{cc}
c_{5} & c_{6} \\
c_{11} & c_{12}
\end{array}\right) \quad(\bmod 26) .
$$

Problem 4 (6.6.9). Let $a, b, c, d, e, f$ be integers mod 26. Consider the following combination of the Hill and affine ciphers: Represent a block of plaintext as a pair $(x, y) \bmod 26$. The corresponding ciphertext $(u, v)$ is

$$
\left(\begin{array}{ll}
x & y
\end{array}\right)\left(\begin{array}{ll}
a & b \\
c & d
\end{array}\right)+\left(\begin{array}{ll}
e & f
\end{array}\right) \equiv\left(\begin{array}{ll}
u & v
\end{array}\right) \quad(\bmod 26) .
$$

Describe how to carry out a chosen plaintext attack on this system (with the goal of finding the key $a, b, c, d, e, f)$. You should state explicitly what plaintext you choose and how to recover the key.

Solution. Use aabaab $=001001$ as the plaintext. Then

$$
\begin{aligned}
& \left(\begin{array}{ll}
0 & 0
\end{array}\right)\left(\begin{array}{ll}
a & b \\
c & d
\end{array}\right)+\left(\begin{array}{ll}
e & f
\end{array}\right) \equiv\left(\begin{array}{ll}
e & f
\end{array}\right) \\
& \left(\begin{array}{ll}
1 & 0
\end{array}\right)\left(\begin{array}{ll}
a & b \\
c & d
\end{array}\right)+\left(\begin{array}{ll}
e & f
\end{array}\right) \equiv\left(\begin{array}{ll}
a+e & b+f
\end{array}\right) \\
& \left(\begin{array}{ll}
0 & 1
\end{array}\right)\left(\begin{array}{ll}
a & b \\
c & d
\end{array}\right)+\left(\begin{array}{ll}
e & f
\end{array}\right) \equiv\left(\begin{array}{ll}
c+e & d+f
\end{array}\right) \quad(\bmod 26) .
\end{aligned}
$$

Now suppose that the ciphertext is $x_{1} x_{2} x_{3} x_{4} x_{5} x_{6}$, where each $x_{i}$ is represented as an integer $\bmod 26$. Then from above, we can get

$$
(a, b, c, d, e, f) \equiv\left(x_{3}-x_{1}, x_{4}-x_{2}, x_{5}-x_{1}, x_{6}-x_{2}, x_{1}, x_{2}\right) \quad(\bmod 26)
$$

Problem 5 (4.6.3). Suppose a message $m$ is chosen randomly from the set of all five-letter English words (so all five-letter words in the dictionary have the same probability) and is encrypted to a ciphertext $c$ using an affine cipher mod 26 , where the key is chosen randomly from the 312 possible keys. Compute the conditional probability

$$
P(m=\text { hello } \mid c=\text { HHGZC }) .
$$

Use the result of this computation to determine whether or not affine ciphers have perfect secrecy.

Solution. Since two l's in hello cannot encrypt to G and Z using the same affine cipher, the conditional probability is zero. Since $P(m=$ hello $)$ is nonzero, this shows that affine ciphers does not have perfect secrecy.

Problem 6 (4.6.13). Alice encrypts the messages $m_{1}$ and $m_{2}$ with the same one-time pad using only capital letters and spaces. (Note. For this problem, assume that each letter or space is converted to a 7 -bit block without the parity bit, according to the ASCII code.) Eve knows this, intercepts the ciphertexts $C_{1}$ and $C_{2}$, and also learns that the decryption of $m_{1}$ is the following.

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THE LETTER * ON THE MAP GIVES THE LOCATION OF THE TREASURE
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Unfortunately for Eve, she cannot read the missing letter $*$. However, the 12 th group of seven bits in $C_{1}$ is 1001101 and the 12th group in $C_{2}$ is 0110101 . Find the missing letter. (Use $\mathrm{A}=1000001, \mathrm{~B}=1000010, \ldots, \mathrm{Z}=1011010$ and space $=0100000$.)

Solution. The 12 th group of seven bits in $m_{1} \oplus m_{2}$ is equal to that in $C_{1} \oplus C_{2}$, which is

$$
1001101 \oplus 0110101=1111000
$$

Since it begins with 1 and the 12 th block of $m_{1}$ is known to be a letter, we can conclude that the 12 th block of $m_{2}$ is the space ( 0100000 ) and consequently the 12 th block of $m_{1}$ is

$$
1111000 \oplus 01000000=1011000=\mathrm{X}
$$

